

**Problem 1** (Long-Run Consensus). Consider the DeGroot learning model with  $N$  agents with initial belief vector  $x(0) = (x_1(0), \dots, x_N(0))$  and an  $N \times N$ , non-negative, row stochastic matrix  $T$  such that, for every period  $t$ , we have

$$x(t) = Tx(t-1).$$

(a) Suppose that  $N = 3$  and

$$T = \begin{pmatrix} \frac{3}{5} & \frac{1}{5} & \frac{1}{5} \\ \frac{1}{4} & \frac{1}{2} & \frac{1}{4} \\ 0 & \frac{1}{3} & \frac{2}{3} \end{pmatrix}$$

What properties of this matrix guarantee that, for any initial belief vector  $x(0)$ , the limit belief  $x^* = \lim_{t \rightarrow \infty} x(t)$  is well-defined? Compute  $x^*$  as a function of  $x(0)$ .

*Solution.* The right-stochastic matrix  $T$  is aperiodic and strongly connected, so we know from lecture that there is a unique limiting belief  $x^*$  that depends only on  $x(0)$ . Moreover, it is given by  $s^\top x(0)$  where the weight vector  $s$  solves

$$sT = s \iff s(T - I) = 0$$

and also satisfies  $\sum_i s_i = 1$ . Solving this linear system of equations by hand, or plugging them into Wolfram Alpha (or a similar tool) gives us  $s = (\frac{5}{22}, \frac{8}{22}, \frac{9}{22})$ .

(b) Suppose that  $N = 6$  and

$$T = \begin{pmatrix} \frac{1}{2} & \frac{1}{6} & \frac{1}{3} & 0 & 0 & 0 \\ \frac{3}{5} & \frac{2}{5} & 0 & 0 & 0 & 0 \\ \frac{3}{11} & \frac{4}{11} & \frac{4}{11} & 0 & 0 & 0 \\ 0 & 0 & 0 & \frac{1}{3} & 0 & \frac{2}{3} \\ 0 & 0 & 0 & \frac{4}{13} & \frac{3}{13} & \frac{6}{13} \\ 0 & 0 & 0 & \frac{4}{7} & 0 & \frac{3}{7} \end{pmatrix}.$$

Without doing any computations, does  $x^* = \lim_{t \rightarrow \infty} x(t)$  exist? Why or why not? If so, which components of the vector  $x^*$  will be identical, and which components may differ?

*Solution.*  $x^*$  exists because the network can be broken into two strongly connected components, and the updating matrix is aperiodic. The first three components of  $x^*$  will be identical, as will the last three components, but typically the first three components will be different from the last three components.

- (c) Prove that, for any  $N$ , if there exists an agent  $i$  such that  $T_{ii} = 1$  and  $T_{ji} > 0$  for all  $j \neq i$ , then  $x_j^* \equiv \lim_{t \rightarrow \infty} x_j(t)$  is well-defined and equal to  $x_i(0)$  for all  $j \neq i$ .

[Hint: Let  $\Delta(t) = \max_{j \in N} |x_i(t) - x_j(t)|$  and let  $\underline{T} = \min_{j \neq i} T_{ji}$ . Prove that  $\Delta(t+1) \leq (1 - \underline{T}) \Delta(t)$  for all  $t$ . Show that this implies that each  $x_j(t)$  must converge to  $x_i(0)$  as  $t \rightarrow \infty$ .]

*Solution.* As suggested by the hint, let's define  $\underline{T} = \min_k T_{ki} > 0$  and  $\Delta(t) = \max_k |x_k(t) - x_i(t)|$ . Firstly, notice that since  $T_{ii} = 1$  and the rows of  $T$  sum to 1,  $T_{ij} = 0$  for  $j \neq i$  and we must have  $x_i(t) = x_i(0)$  for all  $t$  by matrix multiplication. Then we can compute

$$\begin{aligned} |x_i(t+1) - x_j(t+1)| &= |x_i(t) - x_j(t+1)| \\ &= |x_i(t) - (Tx(t))_j| \\ &= x_i(t) - \sum_{k=1}^n T_{jk} x_k(t) \end{aligned}$$

Since  $\sum_{k=1}^n T_{jk} = 1$  this can be rewritten as

$$= \sum_{k=1}^n T_{jk} (x_i(t) - x_k(t))$$

By the triangle inequality  $|a+b| \leq |a| + |b|$ , this is at most

$$\leq \sum_{k=1}^n T_{jk} |x_i(t) - x_k(t)|$$

Since  $\sum_{k \neq i} T_{jk} = 1 - T_{ji} \leq 1 - \underline{T}$  and  $\Delta(t) = \max_k |x_k(t) - x_i(t)|$ , this is at most

$$\leq (1 - \underline{T}) \Delta(t).$$

Since the above holds for every  $j \neq i$ , we can deduce that  $\Delta(t+1) \leq (1 - \underline{T}) \Delta(t)$  which means that, since  $(1 - \underline{T}) < 1$ , we must have  $\Delta(t) \downarrow 0$  as needed. We conclude that for each  $k$ ,  $x_k(t) \rightarrow x_i(0)$  as  $t \rightarrow \infty$ .

**Problem 2** (Clustering). Consider the Erdős-Renyi model with  $n > 1$  nodes and link probability  $p(n)$ , which changes as we add nodes. Let  $p(n) = \lambda/n$  for all  $n$ . Observe that expected degree is held fixed at  $\lambda$ .

- (a) Show that as  $n \rightarrow \infty$  the expected number of triangles in the network converges to  $\frac{1}{6}\lambda^3$ . (Recall that a *triangle* is a triple of nodes  $(i, j, k)$  such that  $g_{ij} = g_{ik} = g_{jk} = 1$ .) Thus, the expected number of triangles hardly depends on  $n$  (once  $n$  is large). Explain how this is possible.

*Solution.* There are  $\binom{n}{3}$  possible triangles.  $\mathbb{P}(g_{ij} = 1, g_{jk} = 1, g_{ik} = 1) = \mathbb{P}(g_{ij} = 1)\mathbb{P}(g_{jk} = 1)\mathbb{P}(g_{ik} = 1) = \left(\frac{\lambda}{n}\right)^3$

$$\binom{n}{3} \frac{\lambda^3}{n^3} = \frac{n(n-1)(n-2)}{6} \left(\frac{\lambda}{n}\right)^3 \rightarrow \frac{\lambda^3}{6} \quad \text{as } n \rightarrow \infty.$$

This is constant because the number of triangles is  $O(n^3)$  and the probability of all three edges occurring in any triangle is  $O\left(\frac{1}{n^3}\right)$ .

- (b) Show that for large  $n$  the expected number of connected triples in the network is approximately  $\frac{1}{2}n\lambda^2$ . (Recall that a *connected triple* is a triple of nodes  $(i, j, k)$  such that  $g_{ij} = g_{ik} = 1$ .)

*Solution.* There are  $\binom{n}{3}$  possible triples. For any  $i, j, k$  let  $I_{ijk}$  be an indicator random variable that is 1 if  $i, j, k$  are a connected triple and 0 otherwise. Note  $I_{ijk}$  takes value 1 with probability  $\binom{3}{2} \left(\frac{\lambda}{n}\right)^2$ , where  $\binom{3}{2}$  comes from needing two of the three edges to be present. So now for the entire graph we have

$$\mathbb{E} \left[ \sum_{i,j,k} I_{ijk} \right] = \sum_{i,j,k} \mathbb{E}[I_{ijk}] = \binom{n}{3} \binom{3}{2} \left(\frac{\lambda}{n}\right)^2 \approx \frac{1}{2}n\lambda^2$$

for large  $n$ .

- (c) Define the *clustering coefficient* for a random network to be the probability that two neighbors of a node are also neighbors of each other. Compute the clustering coefficient for the Erdős-Renyi model with  $p(n) = \lambda/n$ .

*Solution.* In the ER model edges are independently realized so the clustering coefficient is  $\frac{\lambda}{n}$ .

**Problem 3** (Phase Transition). Consider again the Erdős-Renyi model with  $n > 1$  nodes and link probability  $p(n)$ . Let  $A$  denote the event that node 1 has at least  $l \in \mathbb{N}$  neighbors. Show that there is a phase transition for this event with the threshold function  $t(n) = \lambda/n$  for some  $\lambda > 0$ . [Hint: You may need to use the fact that  $(1 + \frac{x}{n})^n \approx \exp(x)$  when  $n$  is large for any  $x \in \mathbb{R}$ .]

*Bonus:* Using a computer language of your choice, code a program that simulates Erdős-Renyi graphs with  $n$  nodes and connection probability  $p$ . Use a simulation with this program to illustrate the phase transition as  $p(n)$  crosses the threshold  $t(n) = 1/n$ . If you “inspect” the networks produced, what other properties do you notice on either side of the threshold?

*Solution.* We need to prove that for any  $l > 0$ :

- (i)  $\mathbb{P}(A|p(n)) \rightarrow 0$  if  $\frac{p(n)}{t(n)} \rightarrow 0$ .
- (ii)  $\mathbb{P}(A|p(n)) \rightarrow 1$  if  $\frac{p(n)}{t(n)} \rightarrow \infty$ .

First assume that  $\frac{p(n)}{t(n)} \rightarrow 0$ . Denote the degree of node 1 by  $d_1$ . Since  $\frac{p(n)}{t(n)} \rightarrow 0$ , the expected degree satisfies

$$\mathbb{E}[d_1] = (n-1)p(n) = \frac{p(n)}{t(n)}t(n)(n-1) \approx \frac{p(n)}{t(n)} \frac{r(n-1)}{n}.$$

Therefore,  $\mathbb{E}[d_1] \rightarrow 0$ . This implies that  $\mathbb{P}(A|p(n)) \rightarrow 0$ , since otherwise the expected degree would be strictly positive.

Next assume that  $\frac{p(n)}{t(n)} \rightarrow \infty$ . It follows that  $p(n) > \frac{r}{n}$  for sufficiently large  $n$ . The probability that  $A$  does not occur can be bounded as follows:

$$\begin{aligned} \mathbb{P}(A^c|p(n)) &= \sum_{k=0}^{l-1} \mathbb{P}(d_1 = k|p(n)) = \sum_{k=0}^{l-1} p(n)^k (1-p(n))^{n-1-k} \binom{n-1}{k} \\ &\leq \sum_{k=0}^{l-1} t(n)^k (1-t(n))^{n-1-k} \binom{n-1}{k} \\ &\leq \sum_{k=0}^{l-1} t(n)^k (1-t(n))^{n-1-k} \frac{n^k}{k!} = \sum_{k=0}^{l-1} \left(\frac{r}{n}\right)^k \left(1 - \frac{r}{n}\right)^{n-1-k} \frac{n^k}{k!} \\ &\approx \sum_{k=0}^{l-1} \exp(-r) \frac{r^k}{k!}. \end{aligned}$$

The second line follows because if the graph was generated using  $t(n)$  instead of  $p(n)$ , each link would be present with smaller probability and hence the probability

that node 1 has less than  $l$  neighbors (the event  $A^c$ ) would be larger. Since the above equation is true for any  $r \in \mathbb{R}^+$ , considering arbitrarily large  $r$ , it follows that  $\mathbb{P}(A^c|p(n)) \rightarrow 0$ , or equivalently that  $\mathbb{P}(A|p(n)) \rightarrow 1$ .

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